

Interactive Mapping Specification with Exemplar Tuples



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1. Motivations and goals

We propose an Interactive Mapping Specification process targeting nonexpert users that:

- bootstraps with a set of exemplar tuples $(E_{\mathbf{S}}, E_{\mathbf{T}})$, corresponding to a limited number of tuples provided as input;
- challenges the non-expert users with simple boolean questions, which are affordable for such users;
- \bullet is guaranteed to always produce a GLAV mapping \mathcal{M}' that general-

4. Atom refinement

Principle:

Remove superfluous atoms in the left-hand sides of normalized tgds.

Exploration space:

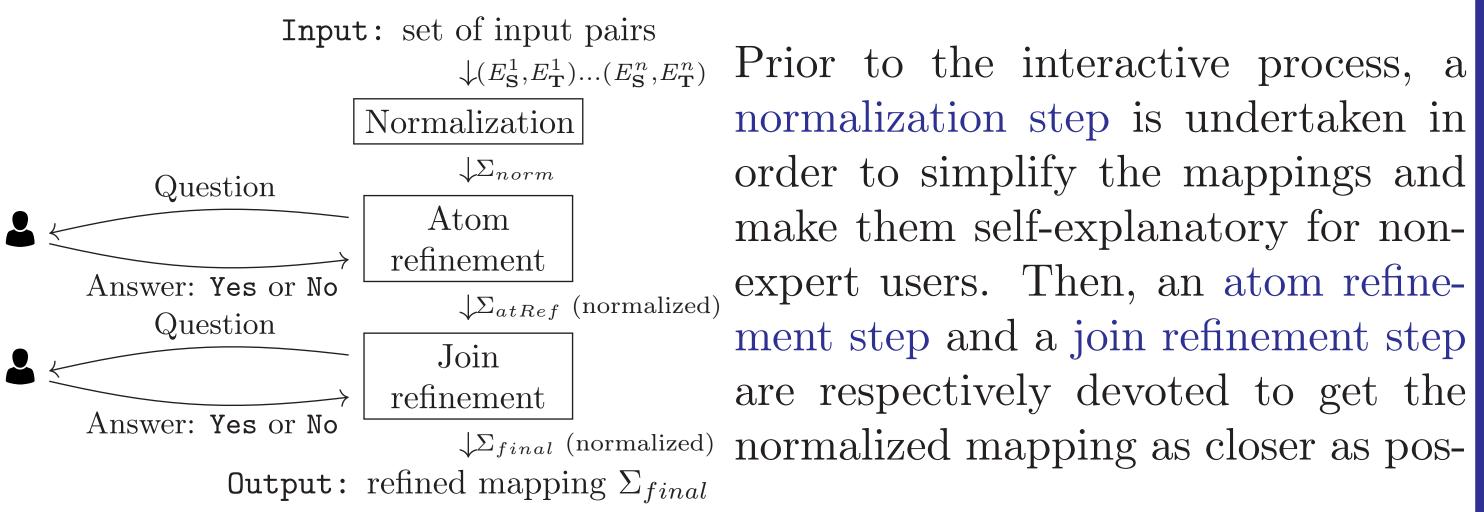
Given a tgd, the possible sets of left-hand side atoms can be represented by a semi-lattice where the lower levels are the smallest valid sets of left-hand side atoms.

User feedback:

During semi-lattice exploration, the user is asked about the target tuples generated using only the subset of source tuples corresponding to the atoms in one node at a time.

izes a mapping \mathcal{M} in the user's mind, which is unknown beforehand.

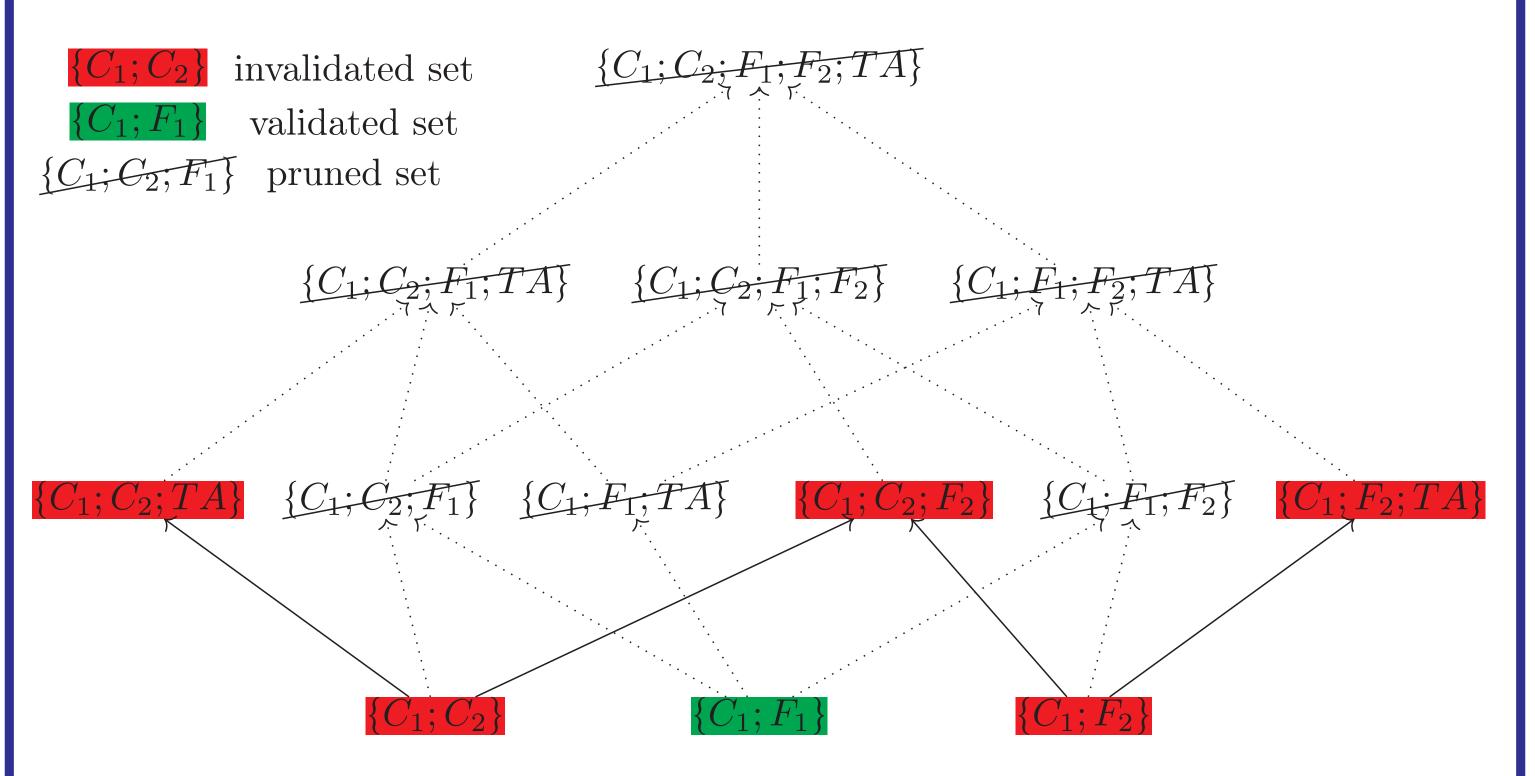
2. Proposed framework



sible to the mapping that the user has in mind. This interactive process leverages simple boolean questions on even smaller refinement-driven tuples, derived from the initial exemplar tuples provided as input. We guarantee that (i) the obtained refined mappings are in normal form and that (ii) they are more general than the canonical mapping.

3. Specifying mappings using user exemplar tuples User exemplar tuples:

Example with a bottom-up breadth-first approach:



As the sole set validated is $\{C_1; F_1\}$, the following tgd is generated: $Company(c1, aa, paris) \land Flight(lyon, paris, c1)$

 $\rightarrow \exists id1, Firm(id1, aa, paris) \land Departure(lyon, id1) \land Arrival(paris, id1)$

Company				\mathbf{Flight}							
$E_{\mathbf{S}}$:	IdCompany		Name	Tou	vn	Departure		Arriva	$l \mid IdCo$	IdCompany	
	'C1'		'AA'	'Par	ris'	'Lyon'		'Paris	' '('C1'	
	'C2'		'Ev'	Lyc	on'	'Paris'		Lyon	' '(C2'	
Travel Agency											
			IdAge	IdAgency		ne Town					
			'A	'A1'		'TC' ']					
	T .•										
$E_{\mathbf{T}}$:	Firm				Departure			Arrival			
	<i>Id</i> 'Id1' 'Id2' 'Id3'	<i>Name</i> 'AA' 'Ev' 'TC'		·is' on'	To	wn	IdFirm	\imath	Town	IdFirm	
					'Ly	ron'	'Id1'		'Paris'	'Id1'	
					'Pa	ris'	'Id2'		'Lyon'	'Id2'	
	100			7.							

Canonical mapping:

 $\mathbf{m}: Company(c1, aa, paris) \land Company(c2, ev, lyon) \land TravelAgency(a1, tc, la)$

- $\land Flight(lyon, paris, c1) \land Flight(paris, lyon, c2)$
- $\rightarrow \exists id1, id2, id3, Firm(id1, aa, paris) \land Departure(lyon, id1) \land Arrival(paris, id1)$
 - $\land Firm(id2, ev, lyon) \land Departure(paris, id2) \land Arrival(lyon, id2) \land Firm(id3, tc, la)$

Normalized mapping:

Before initiating the refinement steps, the canonical mapping is normalized [1] in order to separate unrelated connected components in the right-hand sides. This leads to the following split-reduced mapping:

 $\mathbf{m_a}: Company(c1, aa, paris) \land Company(c2, ev, lyon) \land TravelAgency(a1, tc, la) \land Flight(lyon, paris, c1) \land Flight(paris, lyon, c2)$

5. Join refinement

Principle:

For each tgd generated by atom refinement, identify redundant joins entailed by multiple occurrences of a given variable.

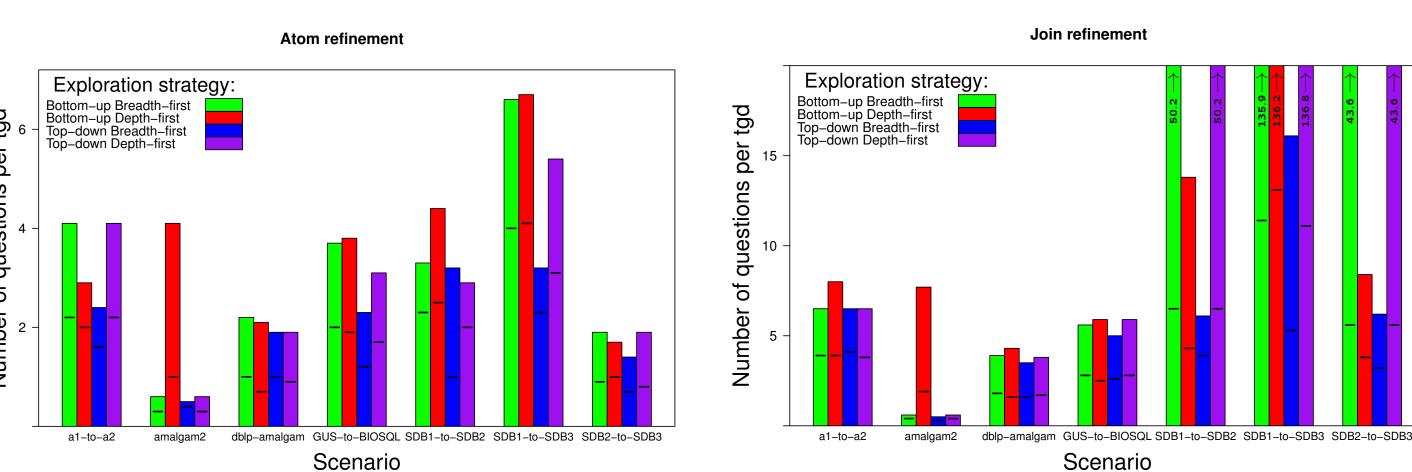
Exploration space:

As with atom refinement, we can build a semi-lattice of partitions representing possible joins between variable occurrences.

User feedback:

Similarly to atom refinement, the user is asked about the validity of small sets of tuples.

6. Experimental study



 $\rightarrow \exists id1, Firm(id1, aa, paris) \land Departure(lyon, id1) \land Arrival(paris, id1)$

- $\mathbf{m_b}: Company(c1, aa, paris) \land Company(c2, ev, lyon) \land TravelAgency(a1, tc, la) \\ \land Flight(lyon, paris, c1) \land Flight(paris, lyon, c2)$
 - $\rightarrow \exists id2, Firm(id2, ev, lyon) \land Departure(paris, id2) \land Arrival(lyon, id2)$
- $\mathbf{m_c}: Company(c1, aa, paris) \land Company(c2, ev, lyon) \land TravelAgency(a1, tc, la) \\ \land Flight(lyon, paris, c1) \land Flight(paris, lyon, c2) \\ \rightarrow \exists id3, Firm(id3, tc, la)$

During normalization, split reduction is followed by the deletion of logically equivalent tgds. In the above example, this leads to the removal of the tgd $\mathbf{m}_{\mathbf{a}}$ (or, equivalently, $\mathbf{m}_{\mathbf{b}}$) in order to obtain Σ_{norm} .

Final mapping after interactive specification:

 $\mathbf{m}_1: Company(c1, aa, paris_1) \land Flight(lyon, paris_2, c1)$

- $\rightarrow \exists id1, Firm(id1, aa, paris_1) \land Departure(lyon, id1) \land Arrival(paris_2, id1)$
- $\mathbf{m}_2: TravelAgency(a1, tc, la) \rightarrow \exists id3, Firm(id3, tc, la)$

- We have used seven real data integration scenarios of the iBench benchmark [2].
- We have simulated the user's ambiguities by adding up to ten superfluous atoms or redundant joins to the initial exemplar tuples, derived from the above scenarios.
- We have tested four exploration strategies : bottom-up and topdown, each one with depth-first and breadth-first variations.

References

- [1] G. Gottlob, R. Pichler, and V. Savenkov:
 - Normalization and optimization of schema mappings. VLDB J., 20(2):277–302, 2011.
- [2] P.C. Arocena, B. Glavic, R. Ciucanu, and R.J. Miller:

The ibench integration metadata generator. Proceedings of VLDB, 9(3):108-119, 2015.